

# **Overview of the x86-64 kernel**

**Andi Kleen, SUSE Labs, Novell  
ak@suse.de**

**Linux Bangalore 2004**

# What's wrong?

---

- x86-64, x86\_64
- AMD64
- EM64T
- IA32e
- IA64
- x64, CT

# Names

- x86-64, x86\_64
- AMD64
- EM64T
- IA32e
- x64
- CT

# **Basics**

---

- 64bit extended x86 architecture
- Can be used with 32bit OS too
  - But 64bit OS is better
- Originally from AMD
- Shipping by AMD and Intel
  - Servers and desktops and even laptops
- Announced by Transmeta and VIA

# **History of the Linux port**

---

- SUSE Labs project
- Started on simulators in 2000
  - Fork from i386
- Was running on early silicon by AMD
- First betas in 2002
- Shipping product (SLES8) in 2003
- Merged into 2.4 in 2002

# **Long mode**

---

- 64bit addressing support
- 64bit instructions
- 8 more integer and SSE2 registers
  - eax -> rax
  - r8-r15, xmm8-xmm15
- RIP relative addressing mode
  - Faster shared libraries
- Compat mode to run 32bit
  - Practically no performance penalty compared to 32bit OS

# An oops

---

```
general protection fault: 0000 [1]
CPU 0
Modules linked in: ....
Pid: 7026, comm: insmod Tainted:
RIP: 0010:[<ffffffffffa073a000>] <ffffffffffa073a000>{:toops3:f2+0}
RSP: 0000:00001000fc79f40 EFLAGS: 00010216
RAX: ffffffa073a010 RBX: fffffff803c4da0 RCX: 000000000101000
RDX: 0000000000000000 RSI: feedbabedeadbeef RDI: feedbabedeadbeef
RBP: ffffffa073a500 R08: 00000100018af010 R09: 000001001ff6d560
R10: 000001001ff6d570 R11: 0000000000000000 R12: fffffff803c4cc0
R13: fffffff803c4cc0 R14: 000000000000000f R15: fffffff8013cb00
FS: 0000002a9588f4c0(0000) GS:fffffff804c6480(0000) knlGS:0000000000000000
CS: 0010 DS: 0000 ES: 0000 CR0: 00000008005003b
CR2: 000000000051b000 CR3: 00000000000101000 CR4: 00000000000006e0
Process insmod (pid: 7026, threadinfo 000001000fc78000, task 000001001d7610b0)
Stack: ffffffa073a019 000001001d7610b0 fffffff80110e47 fffffff8013cb00
    000000000000000f fffffff803c4cc0 fffffff803c4cc0 ffffffa073a500
    fffffff803c4da0 0000000000000000
Call Trace:<ffffffffffa073a019>{:toops3:crash+9} <ffffffffff80110e47>{:child_rip+8}
    <ffffffffff8013cb00>{:msleep+0} <fffffffffa073a010>{:toops3:crash+0}
    <ffffffff80110e3f>{:child_rip+0}

Code: c6 07 01 c3 66 66 66 90 66 66 66 90 48 83 ec 08
RIP <ffffffffffa073a000>{:toops3:f2+0} RSP <000001000fc79f40>
done

0:      c6 07 01          movb    $0x1,(%rdi)
3:      c3                  retq
```

# Some myths

---

- 64bit is bigger
  - Depends on what CPU you optimize for
  - Normally <~10% difference
  - Sometimes code is even smaller
- 64bit is slower
  - Additional registers
  - New modern ABI
  - SSE2
- I don't need 64bit, I have less than 4GB of RAM
  - 32bit limit in practice around 2GB
  - Virtual address space fragments (e.g. thread stacks)
  - IO memory hole needs physical space below 4GB

# **Basics**

---

- Started as a copy of arch/i386, include/asm-i386
- Low level assembly code rewritten
- Code heavily changed for 64bit
  - And only support modern chipsets
- Lots of old cruft removed
  - Workarounds for old hardware bugs
  - No DMI checks so far
  - No APM, no vm86, ...
- Some code shared: MTRR, cpufreq, swiotlb, ...

# **New features**

---

- **NUMA**
  - Based on generic NUMA infrastructure in VM
  - Originally for Opteron only, now also supports ACPI SRAT
  - NUMA API
- **32bit emulation**
  - Based on code from other 64bit ports
- **IOMMU**
- **4level page tables**
  - Before that 512GB limit per process
- **Redesigned machine check handling**

## **Current state**

---

- Widely used
- 2.4 in maintenance mode
- 2.6 production and development

# Porting: basics

---

- Code must be 64bit clean
- long is 64bit now, int stays 32bit
- Pointers in long, not int
  - different from WIN64
- -Wall cleanliness is a good start

# Porting in userspace: /lib64

---

- All 64bit libraries are in lib64
  - 32bit stays in lib
  - Special compat packages for old libraries
- Makefiles often need to be fixed
  - configure --enable-lib-suffix=64
- Not perfect: no include64, bin64
  - Best to have separate library RPMS
  - RPM versions should match

# Porting: IOMMU basics

---

- Some devices cannot address all memory
  - Kick your hardware people if it happens with new hardware
- Driver must map buffers before passing them to hardware
  - Replaces \_\_va, virt\_to\_bus
  - And free them of course
  - Should be used always
- Explicit cache flushing
- Only works for devices with at least 4GB address space
  - Smaller ones need pci\_alloc\_consistent()

# Porting: IOMMU implementation on x86-64

---

- AMD AGP GART IOMMU

- Not a real IOMMU...
- Uses AGP GART functionality in the CPU northbridge
- Reuses half of the AGP aperture by default
- Size depends on BIOS or can be mapped over memory

- Slower swiotlb on Intel

- And some buggy AMD chipsets
- Does memory copies
- Slow

- Remap space is limited

- Sometimes only 64MB
- Can be tuned with kernel command line options and in BIOS
- Best to limit yourself and handle overflows

# Porting: IOMMU functions

---

- pci\_set\_dma\_mask
- pci\_alloc\_consistent for IO memory
  - pci\_free\_consistent
- pci\_map\_sg/pci\_map\_single for dynamic mappings
  - Need 4GB dma mask or better
- pci\_dma\_sync\_{single,sg}\_for\_{device,cpu}

# Porting: IOMMU notes

---

- Check and handle errors
  - Especially in block drivers!
  - pci\_map\_sg returns 0 on error
  - pci\_dma\_mapping\_error for pci\_map\_single
- dma\_\* can be used too for generic bus support
  - pci\_alloc\_consistent -> dma\_alloc\_coherent
  - pci\_map\_single -> dma\_map\_single
  - pci\_map\_sg -> dma\_map\_sg
  - pci\_dma\_mapping\_error -> dma\_mapping\_error
- Documentation/DMA-mapping.txt

# Porting: 32bit emulation basics

---

- 32bit has separate libraries in user space
- 32bit and 64bit always run in different processes
- Kernel has a 32bit emulation layer
- Kernel converts all system calls
  - {fs,net,kernel}/compat.c
- ioctls in drivers need special conversion
- Avoid message passing over read/write

# Porting: 32bit ioctl handler

---

- Needed for x86\_64, ppc64, s390x, ia64, mips64, parisc64
- Kernel does it centrally for most of its own ioctls
  - fs/compat-ioctl.{c,h}
- Drivers can register own ioctl handler
  - register\_ioctl32\_conversion
- Passed through if compatible or converted
- Conversion of structures on user stack
  - Converted from 64bit to compat\_\* types
  - Access using normal \*\_user functions

# What needs conversion?

- long
- pointers
- long long / u64 without natural alignment
  - Different from RISC ports!
- Some fundamental types
  - dev\_t, inode\_t, time\_t, ...

# ioctl conversion functions

- #include <linux/compat.h>
- register\_ioctl32\_conversion()
  - Need unique number
  - Use \_IO\* macros to define ioctls
- copy\_in\_user()
- compat\_alloc\_userspace()
- sys\_ioctl()
- compat\_ptr()

# 32bit conversion example

---

```
#include <linux/compat.h>

struct ppp_idle32 {
    compat_time_t xmit_idle;
    compat_time_t recv_idle;
};

#define PPPIOCGIDLE32           _IOR('t', 63, struct ppp_idle32)

static int ppp_gidle(unsigned int fd, unsigned int cmd, unsigned long arg)
{
    struct ppp_idle __user *idle;
    struct ppp_idle32 __user *idle32;
    __kernel_time_t xmit, recv;
    int err;

    idle = compat_alloc_user_space(sizeof(*idle));
    idle32 = compat_ptr(arg);

    err = sys_ioctl(fd, PPPIOCGIDLE, (unsigned long) idle);

    if (!err) {
        if (get_user(xmit, &idle->xmit_idle) ||
            get_user(recv, &idle->recv_idle) ||
            put_user(xmit, &idle32->xmit_idle) ||
            put_user(recv, &idle32->recv_idle))
            err = -EFAULT;
    }
    return err;
}
```

# **References**

---

- /usr/src/linux/arch/x86\_64, include/asm-x86\_64/
- <http://www.x86-64.org>
- /usr/src/linux/Documentation/DMA-mapping.txt
- discuss@x86-64.org
- Questions?

# Backups

---

# Porting issues: 32bit code with 64bit apps

---

- Direct linking not possible
- All conversion is in the kernel
- Recommended method: several processes, RPC
- Make sure your RPC encoding doesn't assume wordsize
- Example: Konqueror using 32bit plugins with DCOP